# **Introduction to Message Passing Interface (MPI) Programming**

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# **Message Passing Paradigm**

Distributed memory architecture:

network

Each process(or) can only access its dedicated address space. No global shared address space Data exchange and communication between processes is done by explicitly passing messages through a communication



Message passing library:

- Should be flexible, efficient and portable
- Hide communication hardware and software layers from application programmer



# **Message Passing Paradigm**

- Widely accepted standard in HPC / numerical simulation: Message Passing Interface (MPI)
- **Process based approach: All variables are local!**
- Data exchange between processes(a.k.a. tasks): Send/receive messages via MPI library calls
- § This is usually the most tedious but also the most flexible way of parallelization
- MPI is standard for explicit message passing today.



# **A Modern MPI Standard**

- § MPI forum defines MPI standard / library subroutine interfaces
- Beginning: April 1992 Before: vendor specific libraries
- § Latest standard: MPI 3.1 (2015) MPI 4.0 under development
- Members (approx. 60) of MPI standard forum
	- Application programmers
	- Research institutes & Computing centres
	- Manufacturers of supercomputers & software designers
- § Successful free implementation: MPICH, OpenMPI + many others + vendor libraries (Intel, IBM, CRAY)
- § All documents and more pointers available at: www.forum.org/
- § MPI defines more than 500 subroutines typically only 10-30 are used



# **Goals and Scope**

#### § PORTABILITY: Architecture and hardware independent code



- § FORTRAN, C & C++ Interface
- § Provides ´well-defined´and ´safe´data transfer
- Enables development of parallel libraries
- § Support heterogeneous environment (e.g. clusters with heterogeneous compute nodes)



# **Software Architecture**



- Operating system view: **• parallel work done by** tasks/processes ■ Programmer's view: Library routines for
	- coordination
	- communication
	- synchronization
- **User's view: MPI execution** environment provides

■ resource allocation (with the support from LSF, SLURM, openPBS, etc.)



# **Parallel Execution**

§ Tasks run throughout program execution: All variables are local



- § Startup phase: MPI
	- launches tasks
- establishes communication context (communicator) among all tasks
- MPI Point-to-point data transfer:
	- usually between pairs of tasks
	- § usually coordinated
	- may be blocking or non-blocking
- § MPI Collective communication:
	- between all tasks or a subgroup of tasks
	- barrier, reductions, scatter/gather
- § Shutdown by MPI



### **MPI Functions**

- MPI consists of hundreds of functions
- Most users will only use a handful small groups of them
- All functions prefixed with MPI
- C functions return integer error
	- MPI\_SUCCESS if no error

Note: MPI functions will be illustrated in C adapted from ORNL training course



# **More Terminology**

#### ■ Communicator

- An object that represents a group of processes (a.k.a. tasks) than can communicate with each other
- □ MPI Comm type
- □ Predefined communicator: MPI\_COMM\_WORLD
- Rank
	- Within a communicator each process is given a unique integer ID
	- Ranks start at 0 and are incremented contiguously
- Size
	- The total number of ranks in a communicator



- int MPI\_Init( int \*argc, char \*\*\*argv )
- § argc
	- Pointer to the number of arguments
- argv
	- Pointer to argument vector
- Initializes MPI environment
- Must be called before any other MPI call



int MPI\_Finalize( void )

- Cleans up MPI environment
- Calling MPI functions after MPI Finalize is undefined



- int MPI\_Comm\_rank(MPI\_Comm comm, int \*rank)
- comm is an MPI communicator
	- **E** Usually MPI COMM WORLD
- § rank
	- will be set to the rank of the calling process in the communicator of comm



int MPI\_Comm\_size(MPI\_Comm comm, int \*size)

- comm is an MPI communicator
	- □ Usually MPI\_COMM\_WORLD
- $\blacksquare$  size
	- will be set to the number of ranks in the communicator comm



```
#include "stdio.h"
#include "mpi.h"
int main(int argc, char **argv)
{
     int rank, size;
     MPI Init(&argc, &argv);
     MPI_Comm_rank(MPI_COMM_WORLD, &rank);
     <code>MPI\bar{\phantom a}Comm\bar{\phantom a}size(MPI\bar{\phantom a}COMM\bar{\phantom a}WORLD, &size);</code>
     printf("Hello from rank %d of %d total \n", rank, size);
     MPI Finalize();
     return 0;
}
```


- § Compile
	- mpicc wrapper used to link in libraries and includes
	- Uses standard C compiler, such as gcc, under the hood
- \$ mpicc hello\_mpi.c –o hello



§ Run

▫ mpirun –n # ./a.out launches # copies of a.out

\$ mpirun –n 3 ./hello Hello from rank 0 of 3 total Hello from rank 2 of 3 total Hello from rank 1 of 3 total



#### § Run

□ mpirun –n # ./a.out launches # copies of a.out

\$ mpirun –n 3 ./hello Hello from rank 0 of 3 total Hello from rank 2 of 3 total Hello from rank 1 of 3 total Note the order





- Many MPI functions require a datatype
- Built in types for all intrinsic C types □ MPI\_INT, MPI\_FLOAT, MPI\_DOUBLE, …



## **Point to Point**

- Point to Point routines
	- Involves two and only two processes
	- One process explicitly initiates send operation
	- One process explicitly initiates receive operation
	- Several send/receive flavors available
		- § Blocking/non-blocking
		- Combined send-receive
	- Basis for more complicated messaging



### **Point to Point**

int MPI\_Send(void \*buf, int count, MPI\_Datatype datatype, int dest, int tag, MPI\_Comm comm)





### **Point to Point**

int MPI Recv(void \*buf, int count, MPI\_Datatype datatype, int source, int tag, MPI\_Comm comm, MPI\_Status status)





# **Point to Point: try 1**

```
#include "mpi.h" int main(int argc, char **argv)
\{int rank, other rank, tag, send buff, recv buff;
    MPI Status status;
    tag = 5;send buff = 10;
    MPI_Init(&argc, &argv);
    MPI Comm rank(MPI COMM WORLD, &rank);
    if(rank==0)other rank = 1;
    else
        other rank = 0;
    MPI Recv(&recv buff, 1, MPI INT, other rank, tag, MPI COMM WORLD, &status);
    MPI Send(&send buff, 1, MPI INT, other rank, tag, MPI COMM WORLD);
    MPI_Finalize();
    return 0;
}
```


# **Point to Point: try 1**

- § Compile
	- \$ mpicc dead.c –o lock
- § Run \$ mpirun –n 2 ./lock
	- … … … deadlock



### **Deadlock**

#### • The problem

- Both processes (a.k.a. tasks) are waiting to receive a message
- Neither process ever sends a message
- Deadlock as both processes wait forever



## **Point to Point: Fix 1**

```
#include "mpi.h"<br>int main(int argc, char **argv)
       int rank, other rank, recv_buff;<br>MPI Status status;
      MPI_Init(&argc, &argv);<br>MPI_Comm_rank(MPI_COMM_WORLD, &rank);
       if(rank==0){<br>
other rank = 1;
       MPI_Recv(&recv_buff, 1, MPI_INT, other_rank, 5, MPI_COMM_WORLD, &status);<br>MPI_Send(&rank, 1, MPI_INT, other_rank, 5, MPI_COMM_WORLD);<br>other_rank = 0;
              MPI_Send(&rank, 1, MPI_INT, other_rank, 5, MPI_COMM_WORLD);<br>MPI_Recv(&recv_buff, 1, MPI_INT, other_rank, 5, MPI_COMM_WORLD, &status);
       MPI_Finalize();
       return 0;
```


- Non blocking Point to Point functions
	- Allow Send and Receive to not block on CPU
		- § Return before buffers are safe to reuse
	- Can be used to prevent deadlock situations
	- Can be used to overlap communication and computation
	- Calls prefixed with "I", because they return **i**mmediately



int MPI Isend(void \*buf, int count, MPI Datatype datatype, int dest, int tag, MPI\_Request request)





int MPI Irecv(void \*buf, int count,

MPI Datatype datatype, int source, int tag, MPI Comm comm, MPI Request request)





int MPI Wait(MPI Request \*request, MPI Status \*status)

- request
	- The request you want to wait to complete
- status
	- Status struct containing information on completed request
- Will block until specified request operation is complete
- MPI Wait, or similar, must be called if request is used



## **Point to Point: Fix 2**

```
#include "mpi.h"<br>int main(int argc, char **argv)
\{ int rank, other rank, recv buff;
       MPI_Request send_req, recv_req;<br>MPI_Status send_stat, recv_stat;
       MPI_Init(&argc, &argv);<br>MPI_Comm_rank(MPI_COMM_WORLD, &rank);<br>if(rank==0) other_rank = 1;
        else other rank = 0;
       MPI_Irecv(&recv_buff, 1, MPI_INT, other_rank, 5, MPI_COMM_WORLD, &recv_req);<br>MPI_Isend(&rank, 1, MPI_INT, other_rank, 5, MPI_COMM_WORLD, &send_req);
       MPI_Wait(&recv_req, &recv_stat);<br>MPI<sup>-</sup>Wait(&send<sup>-</sup>req, &send<sup>-</sup>stat);
       MPI Finalize();
        ret\overline{u}rn \theta;
```


- Collective routines
	- Involves all processes in communicator
	- All processes in communicator **must** participate
	- Serve several purposes
		- § Synchronization
		- Data movement
		- Reductions
	- Several routines originate or terminate at a single process known as the "root"



int MPI\_Barrier( MPI\_Comm comm )

- Blocks process in comm until all process reach it
- Used to synchronize processes in comm



#### Broadcast





int MPI Bcast(void \*buf, int count, MPI\_Datatype datatype, int root, MPI\_Comm comm)

#### **buf**

Initial address of send buffer **count** Number of elements to send **datatype** Datatype of each element in send buffer

#### **root**

Rank of node that will broadcast buf

**comm**

**Communicator** 





**Scatter** 





int MPI Scatter(const void \*sendbuf, int sendcount, MPI Datatype sendtype, void \*recvbuf int recvcount, MPI Datatype recvtype, int root, MPI\_Comm comm)



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defy convention



**Gather** 







int MPI Gather(const void \*sendbuf, int sendcount, MPI Datatype sendtype, void \*recvbuf int recvcount, MPI Datatype recvtype, int root, MPI\_Comm comm)





#### **Reduce**





int MPI Reduce(const void \*sendbuf, void \*recvbuf int count, MPI Datatype datatype, MPI\_Op op, int root, MPI\_Comm comm)

#### **sendbuf**

Initial address of send buffer **recvbuf**

> Buffer to receive reduced result on root rank

#### **count**

Number of elements in sendbuf

#### **datatype**

Datatype of each element in send buffer

#### **op**

Reduce operation (MPI\_MAX, MPI\_MIN, MPI\_SUM, …)

#### **root**

Rank of node that will receive reduced value in recvbuf

#### **comm**

**Communicator** 



## **Collectives: Example**

```
#include "stdio.h" #include "mpi.h"
int main(int argc, char **argv)<br>{<br>int rank, root, bcast data;
     root = 0;
     if(rank == root)<br>bcast_data = 10;
     MPI_Init(&argc, &argv);<br>MPI_Comm_rank(MPI_COMM_WORLD, &rank);
     MPI Bcast(&bcast data, 1, MPI INT, root, MPI COMM WORLD);
     printf("Rank %d has bcast data = %d\n", rank, bcast data);
     MPI_Finalize();
return 0;<br>}
                                                               Bug here???
```


### **Collectives: Example**

- \$ mpicc collectives.c –o coll
- \$ mpirun –n 4 ./coll
- Rank  $\theta$  has bcast data = 10
- Rank 1 has bcast data =  $10$
- Rank 3 has bcast data =  $10$
- Rank 2 has bcast data =  $10$



## **Submit an MPI Job**

- § #!/bin/bash
- $\bullet$  #
- § #SBATCH --job-name=mpi\_test
- #SBATCH --output=output.txt
- $\blacksquare$
- $\blacksquare$  # Number of MPI tasks
- § #SBATCH -n 6
- $\blacksquare$
- # Number of tasks per node
- § #SBATCH --tasks-per-node=2
- § #
- # Runtime of this jobs is less then 12 hours.
- § #SBATCH --time=12:00:00
- # request to run on general nodes not dev nodes
- § #SBATCH --partition=general
- § #
- **•** mpirun -n 6 sleep 120
- $\blacksquare$  # End of submit file

